On Minimal Doubly Resolving Sets of Circulant Graphs

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Abstract: Consider a simple connected undirected graph $G = (V_G, E_G)$, where V_G represents the vertex set and E_G represents the edge set respectively. A subset B of V_G is called a resolving set if for every two distinct vertices x, y of G there is a vertex v in set B such that $d(x, v) \neq d(y, v)$. The resolving set of minimum cardinality is called metric basis of graph G . This minimal cardinality of metric basis is denoted by $\beta(G)$, and is called metric dimension of G . A subset D of V is called doubly resolving set if for every two vertices x, y of G there are two vertices $u, v \in D$ such that $d(u, x) - d(u, y) \neq d(v, x) - d(v, y)$. A doubly resolving set with minimum cardinality is called minimal doubly resolving set. This minimum cardinality is denoted by $\psi(G)$.

Some partial cases for metric dimension of circulant graph $C_n(1, 2, 3)$ for $n \geq 12$ has been discussed in [21]. Afterwards, problem of finding metric dimension for circulant graph $C_n(1, 2, 3)$, $n \geq 12$ has been completely solved by Borchert et al., in [7].

In this paper, we prove that $\psi(C_n(1, 2, 3)) = \beta(C_n(1, 2, 3)) = \begin{cases} 4 & \text{if } n \not\equiv 1 \pmod{6}, \\ 5 & \text{otherwise.} \end{cases}$

Keywords: resolving set, metric dimension, minimal doubly resolving set, circulant graph.

1. Introduction and preliminary results

Slater and Harary introduced the metric dimension problem independently in [1] and in [2] respectively. This problem has been considered and solved completely/partially for many families of graphs. As an example, one can consult from [10-28]. The applications of metric dimension includes network discovery and verification [3], geographic routing protocols and robot navigation [4], connected joints in graphs, and chemistry.

Consider a simple connected undirected graph $G = (V_G, E_G)$, where V_G and E_G denote the set of vertices and set of edges of G , respectively. The distance between vertices x and y is denoted by $d(x, y)$ which is the length of the shortest path between x and y . We say that a vertex v resolve two vertices x and y of G if $d(x, v) \neq d(y, v)$. A subset B of V_G is called a resolving set if every two distinct vertices of graph G are resolved by some vertex in set B . This concept can be explained in another terminology also, which is as follows:

Consider an ordered subset $B = \{x_1, x_2, \dots, x_p\}$ of vertices of G . For an arbitrary vertex y of G , we have the following p -tuple

$$r(y|B) = (d(y, x_1), d(y, x_2), \dots, d(y, x_p))$$

which is called the representation of vertex y or vector of metric coordinates of y with respect to B . The set B is called a resolving set if vector of metric coordinates of each vertex with respect to set B is unique. The resolving set of minimum cardinality is called metric basis of graph G . This minimal cardinality of metric basis is denoted by $\beta(G)$, and is called metric dimension of G .

The notion of doubly resolving set of graph G was introduced by Caceres et al. [10] in the following way. Consider a graph G of order at least 2. Two vertices x, y are said to doubly resolved by vertices x', y' of G if

$$d(x, x') - d(x, y') \neq d(y, x') - d(y, y').$$

An ordered subset $D = \{x_1, x_2, \dots, x_q\}$ of V_G is called a doubly resolving set if every two distinct vertices of G are doubly resolved by some two vertices in set D , i.e., for each pair of vertices $x, y \in V_G$ we have

$$r(x|D) - r(y|D) \neq \lambda I,$$

where λ is an integer and I denotes the unit vector $(1, 1, \dots, 1)$. Minimal doubly resolving set of graph G is a doubly resolving set with minimal cardinality. This minimum cardinality is denoted by $\psi(G)$. Observe that, if vertices x', y' doubly resolve the vertices x, y , then $d(x, x') - d(x, y') \neq 0$ or $d(y, x') - d(y, y') \neq 0$. This shows that x' or y' resolve x, y , which follows that a doubly resolving set is also a resolving set, hence $\beta(G) \leq \psi(G)$. In this way, these sets constitute a useful tool for obtaining upper bounds on the metric dimension of graphs. The metric dimension problem and minimal doubly resolving set problem are NP-hard. The proofs can be found in [6] and [8], respectively. The problem of finding minimal doubly resolving set for different families of graphs has been studied in [5] and [9].

The circulant graphs constitute an important class of graphs which can be used in the design of local area networks [29]. Let n, m and $t_1, t_2, \dots, t_m$ be positive integers, $1 \leq t_i \leq \lfloor n/2 \rfloor$ and $t_i \neq t_j$ for all $1 \leq i < j \leq m$. The circulant graph $C_n(t_1, \dots, t_m)$ can be formed by taking set of vertices $V = \{v_1, \dots, v_n\}$ and set of edges $E = \{v_i v_{i+t_j} \mid 1 \leq i \leq n, 1 \leq j \leq m\}$ with indices taken modulo n . The numbers $t_1, \dots, t_m$ are called

Fig. 1: The circulant graph $C_{12}(1,2,3)$.

the generators and we say that the edge $v_i v_{i+t_j}$ is of type t_j . The graph $C_{12}(1,2,3)$ is shown in Fig. 1 below.

The circulant graph $C_n(t_1, \dots, t_m)$ is a regular graph of degree r , where

$$r = \begin{cases} 2m-1 & \text{if } \frac{n}{2} \in \{t_1, \dots, t_m\}, \\ 2m & \text{otherwise.} \end{cases}$$

In this paper, we explicitly determine the minimal doubly resolving sets for circulant graph $C_n(1,2,3)$ for $n \geq 12$. Also we prove that $\beta(C_n(1,2,3)) = \psi(C_n(1,2,3))$ for $n \geq 12$.

2. The minimal doubly resolving sets for circulant graph $C_n(1,2,3)$ for $n \geq 12$

We have

$$\psi(C_n(1,2,3)) \geq \beta(C_n(1,2,3)) = \begin{cases} 4 & \text{if } n \not\equiv 1 \pmod{6}, \\ 5 & \text{otherwise.} \end{cases}$$

for $n \geq 12$, see [7].

Let us denote the set $S_i(v_j) = \{w \in V(C_n(1,2,3)) \mid d(v_j, w) = i\}$, which is the set of vertices in $V(C_n(1,2,3))$ at distance i from v_j . The Table 1 displays the sets $S_i(v_j)$ for circulant graph $C_n(1,2,3)$, where $n \geq 12$.

Table 1: $S_i(v_j)$ for $C_n(1,2,3)$ for $n \geq 12$.

n	i	$S_i(v_j)$
$6k$ ($k \geq 2$)	0	$\{v_j\}$
	$1 \leq i \leq k-1$	$\{v_{3i-1}, v_{3i}, v_{3i+1}, v_{6k+3-3i}, v_{6k+2-3i}, v_{6k+1-3i}\}$
	k	$\{v_{3k-1}, v_{3k}, v_{3k+1}, v_{3k+2}, v_{3k+3}\}$
$6k+1$ ($k \geq 2$)	0	$\{v_j\}$
	$1 \leq i \leq k$	$\{v_{3i-1}, v_{3i}, v_{3i+1}, v_{6k+4-3i}, v_{6k+3-3i}, v_{6k+2-3i}\}$
	$k+1 \leq i$	$\emptyset$
$6k+2$ ($k \geq 2$)	0	$\{v_j\}$
	$1 \leq i \leq k$	$\{v_{3i-1}, v_{3i}, v_{3i+1}, v_{6k+5-3i}, v_{6k+4-3i}, v_{6k+3-3i}\}$
	$k+1 \leq i$	$\{v_{3k+2}\}$
$6k+3$ ($k \geq 2$)	0	$\{v_j\}$
	$1 \leq i \leq k$	$\{v_{3i-1}, v_{3i}, v_{3i+1}, v_{6k+6-3i}, v_{6k+5-3i}, v_{6k+4-3i}\}$
	$k+1 \leq i$	$\{v_{3k+2}, v_{3k+3}\}$
$6k+4$ ($k \geq 2$)	0	$\{v_j\}$
	$1 \leq i \leq k$	$\{v_{3i-1}, v_{3i}, v_{3i+1}, v_{6k+7-3i}, v_{6k+6-3i}, v_{6k+5-3i}\}$
	$k+1 \leq i$	$\{v_{3k+2}, v_{3k+3}, v_{3k+4}\}$
$6k+5$ ($k \geq 2$)	0	$\{v_j\}$
	$1 \leq i \leq k$	$\{v_{3i-1}, v_{3i}, v_{3i+1}, v_{6k+8-3i}, v_{6k+7-3i}, v_{6k+6-3i}\}$
	$k+1 \leq i$	$\{v_{3k+2}, v_{3k+3}, v_{3k+4}, v_{3k+5}\}$

Theorem 2.1. $\psi(C_n(1,2,3)) = 4$ for $n \equiv 0, 2, 3, 4, 5 \pmod{6}, n \geq 12$.

Proof. We need to show that $\psi(C_n(1,2,3)) \leq 4$ for $n \equiv 0, 2, 3, 4, 5 \pmod{6}, n \geq 12$. So it suffices to find a doubly resolving set of cardinality 4 in each case. Let us first consider the case when $n \equiv 0 \pmod{6}$, i.e., $n = 6k$ for $k \geq 2$. Using the sets $S_i(v_l)$ from Table 1, the Table 2 displays the vectors of metric coordinates of every vertex of $C_n(1,2,3)$ with respect to the set $D^* = \{v_l, v_{3l}, v_{3l+1}, v_{3k+1}\}$.

Table 2: Vectors of metric coordinates for $C_n(1,2,3)$ for $n=6k, k \geq 3$.

i	$S_i(v_l)$	metric coordinates w.r.t. $D^* = \{v_l, v_{3l}, v_{3l+1}, v_{3k+1}\}$
0	v_l	$(0, 1, 2, k)$
1	v_2	$(1, 1, 1, k)$
	v_3	$(1, 0, 1, k)$
	v_4	$(1, 1, 1, k-1)$
	v_{6k}	$(1, 1, 2, k)$
	v_{6k-1}	$(1, 2, 2, k)$
	v_{6k-2}	$(1, 2, 3, k-1)$
$2 \leq i \leq k-2$ ($k \geq 4$)	v_{3i-1}	$(i, i-1, i-2, k+1-i)$
	v_{3i}	$(i, i-1, i-1, k+1-i)$
	v_{3i+1}	$(i, i, i-1, k-i)$
	$v_{6k+3-3i}$	$(i, i, i+1, k+1-i)$
	$v_{6k+2-3i}$	$(i, i+1, i+1, k+1-i)$
	$v_{6k+1-3i}$	$(i, i+1, i+2, k-i)$
$k-1$	v_{3k-4}	$(k-1, k-2, k-3, 2)$
	v_{3k-3}	$(k-1, k-2, k-2, 2)$
	v_{3k-2}	$(k-1, k-1, k-2, 1)$
	v_{3k+6}	$(k-1, k-1, k, 2)$
	v_{3k+5}	$(k-1, k, k, 2)$
	v_{3k+4}	$(k-1, k, k, 1)$
k	v_{3k-1}	$(k, k-1, k-2, 1)$
	v_{3k}	$(k, k-1, k-1, 1)$
	v_{3k+1}	$(k, k, k-1, 0)$
	v_{3k+2}	$(k, k, k-1, 1)$
	v_{3k+3}	$(k, k, k, 1)$

Note that for $k = 2$, it can be checked directly that $\{v_l, v_{3l}, v_{3l+1}, v_{3k+1}\}$ is a doubly resolving set.

In the same way, using Table 1, the Tables 3 - 6 display the vectors of metric coordinates of vertices of $C_n(1,2,3)$ for $n \equiv 2, 3, 4, 5 \pmod{6}$ with respect to the set $D^* = \{v_l, v_{3l}, v_{3l+1}, v_{3k+1}\}$, $D^* = \{v_l, v_{3l}, v_{3l+1}, v_{6k+4}\}$ and $D^* = \{v_l, v_{3l}, v_{3k+1}, v_{3k+7}\}$, respectively.

From Tables 2 to 6, it can be verified directly that if two vertices x, y belongs to $S_i(v_l)$ for some i , then

$$r(x|D^*) - r(y|D^*) \neq 0I,$$

where I denotes the unit vector. Also if $x \in S_i(v_l)$ and $x \in S_j(v_l)$ for $i \neq j$, then

$$r(x|D^*) - r(y|D^*) \neq (i-j)I.$$

Table 3: Vectors of metric coordinates for $C_n(1,2,3)$ for $n=6k+2, k \geq 2$.

i	$S_i(v_l)$	metric coordinates w.r.t. $D^* = \{v_l, v_{3l}, v_{3k+1}, v_{3k+4}\}$
0	v_l	$(0, 1, k, k)$
1	v_2	$(1, 0, k, k)$
	v_3	$(1, 1, k, k+1)$
	v_4	$(1, 1, k-1, k)$
	v_{6k+2}	$(1, 1, k+1, k)$
	v_{6k+1}	$(1, 1, k, k-1)$
	v_{6k}	$(1, 2, k, k-1)$
$2 \leq i \leq k-1$ ($k \geq 3$)	v_{3i-1}	$(i, i-1, k+1-i, k+2-i)$
	v_{3i}	$(i, i, k+1-i, k+2-i)$
	v_{3i+1}	$(i, i, k-i, k+1-i)$
	$v_{6k+5-3i}$	$(i, i, k+2-i, k+1-i)$
	$v_{6k+4-3i}$	$(i, i, k+1-i, k-i)$
	$v_{6k+3-3i}$	$(i, i+1, k+1-i, k-i)$
k	v_{3k-1}	$(k, k-1, 1, 2)$
	v_{3k}	$(k, k, 1, 2)$
	v_{3k+1}	$(k, k, 0, 1)$
	v_{3k+5}	$(k, k, 2, 1)$
	v_{3k+4}	$(k, k, 1, 0)$
	v_{3k+3}	$(k, k+1, 1, 1)$
$k+1$	v_{3k+2}	$(k+1, k, 1, 1)$

Table 4: Vectors of metric coordinates for $C_n(1,2,3)$ for $n=6k+3, k \geq 2$.

i	$S_i(v_l)$	metric coordinates w.r.t. $D^* = \{v_l, v_{3l}, v_{3k+1}, v_{3k+5}\}$
0	v_l	$(0, 1, k, k)$
1	v_2	$(1, 0, k, k)$
	v_3	$(1, 1, k, k+1)$
	v_4	$(1, 1, k-1, k+1)$
	v_{6k+3}	$(1, 1, k+1, k)$
	v_{6k+2}	$(1, 1, k+1, k-1)$
	v_{6k+1}	$(1, 2, k, k-1)$
$2 \leq i \leq k-1$ ($k \geq 3$)	v_{3i-1}	$(i, i-1, k+1-i, k+2-i)$
	v_{3i}	$(i, i, k+1-i, k+2-i)$
	v_{3i+1}	$(i, i, k-i, k+2-i)$
	$v_{6k+6-3i}$	$(i, i, k+2-i, k+1-i)$
	$v_{6k+5-3i}$	$(i, i, k+2-i, k-i)$
	$v_{6k+4-3i}$	$(i, i+1, k+1-i, k-i)$
k	v_{3k-1}	$(k, k-1, 1, 2)$
	v_{3k}	$(k, k, 1, 2)$
	v_{3k+1}	$(k, k, 0, 2)$
	v_{3k+6}	$(k, k, 2, 1)$
	v_{3k+5}	$(k, k, 2, 0)$
	v_{3k+4}	$(k, k+1, 1, 1)$
$k+1$	v_{3k+2}	$(k+1, k, 1, 1)$
	v_{3k+3}	$(k+1, k+1, 1, 1)$

Table 5: Vectors of metric coordinates for $C_n(1,2,3)$ for $n=6k+4$, $k \geq 2$.

i	$S_i(v_i)$	metric coordinates w.r.t. $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{6k+4}\}$
0	v_i	$(0, 1, k, 1)$
$1 \leq i \leq k$	v_{3i-1} v_{3i} v_{3i+1} $v_{6k+7-3i}$ $v_{6k+6-3i}$ $v_{6k+5-3i}$	$(i, i-1, k+1-i, i)$ $(i, i, k+1-i, i)$ $(i, i, k-i, i+1)$ $(i, i, k+2-i, i-1)$ $(i, i, k+2-i, i)$ $(i, i+1, k+2-i, i)$
$k+1$	v_{3k+2} v_{3k+3} v_{3k+4}	$(k+1, k, 1, k+1)$ $(k+1, k+1, 1, k+1)$ $(k+1, k+1, 1, k)$

Table 6: Vectors of metric coordinates for $C_n(1,2,3)$ for $n=6k+5$, $k \geq 2$.

i	$S_i(v_i)$	metric coordinates w.r.t. $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{3k+7}\}$
0	v_i	$(0, 1, k, k)$
1	v_2 v_3 v_4 v_{6k+5} v_{6k+4} v_{6k+3}	$(1, 0, k, k)$ $(1, 1, k, k+1)$ $(1, 1, k-1, k+1)$ $(1, 1, k+1, k)$ $(1, 1, k+1, k-1)$ $(1, 2, k+1, k-1)$
$2 \leq i \leq k-1$ ($k \geq 3$)	v_{3i-1} v_{3i} v_{3i+1} $v_{6k+8-3i}$ $v_{6k+7-3i}$ $v_{6k+6-3i}$	$(i, i-1, k+1-i, k+3-i)$ $(i, i, k+1-i, k+3-i)$ $(i, i, k-i, k+2-i)$ $(i, i, k+3-i, k+1-i)$ $(i, i, k+2-i, k-i)$ $(i, i+1, k+2-i, k-i)$
k	v_{3k-1} v_{3k} v_{3k+1} v_{3k+8} v_{3k+7} v_{3k+6}	$(k, k-1, 1, 3)$ $(k, k, 1, 3)$ $(k, k, 0, 2)$ $(k, k, 3, 1)$ $(k, k, 2, 0)$ $(k, k+1, 2, 1)$
$k+1$	v_{3k+2} v_{3k+3} v_{3k+4} v_{3k+5}	$(k+1, k, 1, 2)$ $(k+1, k+1, 1, 2)$ $(k+1, k+1, 1, 1)$ $(k+1, k+1, 2, 1)$

Thus $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{3k+7}\}$, $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{3k+4}\}$, $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{3k+5}\}$, $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{6k+4}\}$ and $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{3k+7}\}$ are doubly resolving sets (indeed minimal doubly resolving sets) of $C_n(1,2,3)$ for $n \equiv 0, 2, 3, 4, 5 \pmod{6}$ respectively. Hence Theorem 2.1 holds.

Theorem 2.2. $\psi(C_n(1,2,3)) = 5$ for $n \equiv 1 \pmod{6}$, $n \geq 12$.

Proof. The proof is same as the proof of Theorem

2.1. We need to show that $\psi(C_n(1,2,3)) \leq 5$ for $n \equiv 1 \pmod{6}$, $n \geq 12$. For this, consider $n=6k+1$, $k \geq 2$, so the target is to find a doubly resolving set of cardinality 5 for each k . It can be proved by a direct check that $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{6k+2}\}$ and $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{11}\}$ are doubly resolving sets corresponding to $k = 2$ and $k = 3$, respectively. Using the sets $S_i(v_i)$ from Table 1, the Table 7 displays the vectors of metric coordinates of every vertex of $C_n(1,2,3)$, with respect to the set $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{6k+2}\}$.

Table 7: Vectors of metric coordinates for $C_n(1,2,3)$ for $n=6k+1$, $k \geq 4$.

i	$S_i(v_i)$	metric coordinates w.r.t. $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{6k+2}\}$
0	v_i	$(0, 1, 2, 2, k)$
1	v_2 v_3 v_4 v_{6k+1} v_{6k} v_{6k-1}	$(1, 1, 1, 2, k)$ $(1, 0, 1, 1, k)$ $(1, 1, 1, 1, k)$ $(1, 1, 2, 2, k)$ $(1, 2, 2, 3, k)$ $(1, 2, 3, 3, k-1)$
2	v_5 v_6 v_7 v_{6k-2} v_{6k-3} v_{6k-4}	$(2, 1, 0, 1, k-1)$ $(2, 1, 1, 0, k-1)$ $(2, 2, 1, 1, k-1)$ $(2, 2, 3, 3, k-1)$ $(2, 3, 3, 4, k-1)$ $(2, 3, 4, 4, k-2)$
$3 \leq i \leq k-2$ ($k \geq 5$)	v_{3i-1} v_{3i} v_{3i+1} $v_{6k+4-3i}$ $v_{6k+3-3i}$ $v_{6k+2-3i}$	$(i, i-1, i-2, i-2, k+1-i)$ $(i, i-1, i-1, i-2, k+1-i)$ $(i, i, i-1, i-1, k+1-i)$ $(i, i, i+1, i+1, k+1-i)$ $(i, i+1, i+1, i+2, k+1-i)$ $(i, i+1, i+2, i+2, k-i)$
$k-1$	v_{3k-4} v_{3k-3} v_{3k-2} v_{3k+7} v_{3k+6} v_{3k+5}	$(k-1, k-2, k-3, k-3, 2)$ $(k-1, k-2, k-2, k-3, 2)$ $(k-1, k-1, k-2, k-2, 2)$ $(k-1, k-1, k, k, 2)$ $(k-1, k, k, k, 2)$ $(k-1, k, k, k, 1)$
k	v_{3k-1} v_{3k} v_{3k+1} v_{3k+4} v_{3k+3} v_{3k+2}	$(k, k-1, k-2, k-2, 1)$ $(k, k-1, k-1, k-2, 1)$ $(k, k, k-1, k-1, 1)$ $(k, k, k, k, 1)$ $(k, k, k, k-1, 1)$ $(k, k, k-1, k-1, 0)$

From Table 7, it can be seen that the difference between vectors of metric coordinates, of any two chosen vertices, is not an integer multiple of unit vector I . Therefore, the set $D^* = \{v_i, v_{3i}, v_{3k+1}, v_{6k+2}\}$ is doubly resolving set (indeed minimal doubly

resolving set) of the circulant graph $C_n(1,2,3)$ and hence Theorem 2.2 holds.

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